## ETH

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## Discrete Event Systems Exercise $3^{1}$

## 1 Regular Languages and Finite Automaton

Consider the NFA $A$ in Figure 1 and assume that $\Sigma=\{0,1\}$.


Figure 1: NFA $A$.
(i) Transform the NFA into an equivalent deterministic finite automaton.
(ii) Which regular language is accepted by $A$ ?

## 2 Non-Regular Languages

(i) Consider the following language $L_{1}$ :

$$
L_{2}=\left\{0^{a} 1^{b} 0^{c} 1^{d} \mid a, b, c, d \geq 0 \text { and } a=1, b=2, \text { and } c=d\right\} .
$$

Is the language $L_{1}$ regular? Prove your answer!
(ii) Consider the following slightly adapted language $L_{2}$ :

$$
L_{2}=\left\{0^{a} 1^{b} 0^{c} 1^{d} \mid a, b, c, d \geq 0 \text { and if } a=1 \text { and } b=2, \text { then } c=d\right\} .
$$

Is the language $L_{2}$ regular? Be careful when proving your answer!

[^0]
## 3 Adapting a Finite Automaton

Consider the DFA in Figure 3, which accepts the language $L$ and let the alphabet be $\Sigma=\{0,1\}$. Further, let $\Phi(L)$ be defined as $\Phi(L)=\left\{w \in \Sigma^{*}\left|\exists x \in \Sigma^{*},|x|=|w|\right.\right.$ and $\left.w x \in L\right\}$. That is, $\Phi(L)$ denotes the set of first halfs of all strings in $L$.


Figure 2: DFA $B$.
(i) Give a regular expression that describes the language $L$.
(ii) Construct a DFA which accepts a string $w$ if and only if $w \in \Phi(L)$.


[^0]:    ${ }^{1}$ All problems in this series have appeared in previous exams.

