## EHHzürich



## Maximally Expressive GNNs for Outerplanar Graphs

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## GNNs

MPNN: $u_{1+1} \Phi\left(u_{1}, \Psi\left(\operatorname{Nei}_{,}(u)\right)\right)$
Graph Neural Network's goal would be to only consider the isomorphism class of a graph (i.e. 2 isomorphic graphs should give the same output).

1-WL Test

$$
\mathrm{u}_{\mathrm{t}+1} \leftarrow\left(\mathrm{u}_{\mathrm{t}}, \operatorname{Nei} \mathrm{i}_{\mathrm{t}}(\mathrm{u})\right)
$$



## 1-WL Test



1-WL Test


1-WL Test


## 1-WL Test

If $\mathrm{WL}(\mathrm{G}) \neq \mathrm{WL}(\mathrm{H})$, then G and H are not isomorphic However, if $\mathrm{WL}(\mathrm{G})=\mathrm{WL}(\mathrm{H}), \mathrm{G}$ and H might not be isomorphic

$\{4$ blue, 2 red $\}=\{4$ blue, 2 red $\}$

## GNNs

Problem: MPNN's expressiveness is bounded by the 1-WL test. In particular, GINs (Graph Isomorphism Networks) are proven to be as expressive as the 1-WL test.


## GNNs

Can we restrict ourselves to a simpler subclass of graphs?

## Planar graphs

Can be drawn on the plane in such a way that no edges cross each other


Non-planar


## Outerplanar graphs

Planar graph that can be drawn so that no vertex is "trapped" inside the edges of the graph

Outerplanar


Non-outerplanar


## Biconnected graph

Connected graph that is not broken into disconnected pieces by deleting any single vertex

Biconnected


Non-biconnected


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Connected graph that is not broken into disconnected pieces by deleting any single vertex

Biconnected


Non-biconnected


## Hamiltonian cycle

Cycle which goes over each node exactly once


## Hamiltonian cycle

Cycle which goes over each node exactly once
Theorem 1: Biconnected outerplanar graphs have a unique Hamiltonian cycle that can be found in linear time


## Hamiltonian cycle (2 directed variants)



## HALs (Hamiltonian Adjacency Lists)

Annotating each node with the sorted distances $\mathrm{d}_{\mathrm{c}}$ to all its neighbors on the two directed variants of the Hamiltonian cycle C .


$$
d_{c}(e, d)=1
$$

## HALs (Hamiltonian Adjacency Lists)

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$$
d_{c}(e, a)=4
$$

## HALs (Hamiltonian Adjacency Lists)

Annotating each node with the sorted distances $\mathrm{d}_{\mathrm{c}}$ to all its neighbors on the two directed variants of the Hamiltonian cycle C .


$$
\begin{aligned}
& d_{c}(a, e)=1 \\
& d_{c}(a, b)=4
\end{aligned}
$$

## HALs (Hamiltonian Adjacency Lists)

Annotating each node with the sorted distances $\mathrm{d}_{\mathrm{c}}$ to all its neighbors on the two directed variants of the Hamiltonian cycle C .


$$
d_{c}(a, d)=2
$$

## HALs (Hamiltonian Adjacency Lists)

Annotating each node with the sorted distances $\mathrm{d}_{\mathrm{c}}$ to all its neighbors on the two directed variants of the Hamiltonian cycle $C$.


Theorem 2: Two biconnected outerplanar graphs G and H with HAL and reverse sequences $S_{G}, S_{H}$ and $R_{G}, R_{H}$ are isomorphic, iff $S_{G}$ is a cyclic shift of $S_{H}$ or $R_{H}$.


## Idea

Let's build a transformation to make $\mathrm{S}_{\mathrm{G}}$ and $\mathrm{R}_{\mathrm{G}}$ recognizable by WL test

## Agenda

1. Find transformation CAT* that guarantees maximal expressiveness for biconnected outerplanar graphs
2. Extend CAT* to build transformation CAT that covers all outerplanar graphs
3. Use CAT to boost expressiveness of GNNs

## CAT*

Transformation of a biconnected outerplanar graph


## CAT*

1. Find directed Hamiltonian cycle C


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## CAT*

3. Add edges not in the Hamiltonian cycle in both directions


## CAT*

4. Label the new edges ( $u, v$ ) with their distance $d_{c}(v, u)$ according to HAL


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4. Label the new edges ( $u, v$ ) with their distance $d_{c}(v, u)$ according to HAL


## CAT*

2. Give on the directed Hamiltonian cycle C all the edges weight 1


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## CAT*



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## CAT*

Repeat the whole procedure on the reverse hamiltonian cycle $\overleftarrow{\mathrm{C}}$


## CAT*

The output of this transformation is 2 connected components


## CAT*

Theorem 3: Two biconnected outerplanar graphs G and H are isomorphic, iff $\mathrm{WL}\left(\mathrm{CAT}^{*}(\mathrm{G})\right)=\mathrm{WL}\left(\mathrm{CAT}^{*}(\mathrm{H})\right)$


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## Next step:

Let's extend this to all outerplanar graphs.

Biconnected components and blocks
A biconnected component is a maximal biconnected subgraph (with at least 3 nodes)


## Biconnected components and blocks

A biconnected component is a maximal biconnected subgraph (with at least 3 nodes)
A block is a biconnected outerplanar component


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Not maximal


## CAT

Transformation of an outerplanar graph


## CAT

1. Identify blocks


## CAT

2. Let $F$ be the graph induced by the edges that are in none of the blocks


## CAT

2. Let $F$ be the graph induced by the edges that are in none of the blocks

3. For all blocks $B_{i}$


## CAT

3. Let's start with block $B_{1}$


## CAT

3.1 Add the 2 connected components $\mathrm{B}_{\mathrm{i}}, \mathrm{B}_{\mathrm{i}}^{\prime}$ from $\mathrm{CAT}^{*}\left(\mathrm{~B}_{\mathrm{i}}\right)$


## CAT

3.4 For all pairs of nodes $v, \overleftarrow{v}$ in $B_{i}^{\prime}, B_{i}^{\prime}$, add a node $p_{v}$


## CAT

3.4 Add edges edges $\left\{p_{v}, v\right\},\left\{p_{v}, \overleftarrow{v}\right\}$


## CAT

3.5 Add a node $\mathrm{b}_{\mathrm{i}}$


## CAT

3.5 Add edges $\left\{\mathrm{b}_{\mathrm{i}}, \mathrm{p}_{\mathrm{v}}\right\}, \forall \mathrm{v} \in \mathrm{V}\left(\mathrm{B}_{\mathrm{i}}\right)$


## CAT <br> 3.2 Let $A_{i}:=V\left(B_{i}\right) \cap V(F)$



## CAT

3.6 Color in orange all green nodes in A


## CAT

3.6 Connect every node from A to its water green counterpart


## CAT

3. Repeat for every block $B_{i}$


## CAT

4. Add a node g


## CAT

4. Add edges $\left\{\mathrm{g}, \mathrm{b}_{\mathrm{i}}\right\}$ for all nodes $\mathrm{b}_{\mathrm{i}}$


## CAT

Theorem 4: Outerplanar graphs G and H are isomorphic, iff $\mathrm{WL}(\mathrm{CAT}(\mathrm{G}))=$ WL(CAT(H))

## CAT

Linear time complexity:

- Time complexity dominated by computation of blocks and their Hamiltonian cycles (both linear)
- We only add a linear number of nodes and edges. $\mathrm{O}(|\mathrm{V}|+|\mathrm{E}|)$


## Experimental results

Table 4: Predictive performance of MPNNs with and without CAT on different molecular benchmark datasets. Arrows indicate whether smaller $(\downarrow)$ or bigger $(\uparrow)$ results are better. Bold entries are an MPNN with CAT that outperforms the same MPNN without CAT.

| Dataset $\rightarrow$ | ZINC | MOLHIV | MOLBACE | MOLBBBP | MOLSIDER |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\downarrow$ Model | MAE $\downarrow$ | ROC-AUC $\uparrow$ | ROC-AUC $\uparrow$ | ROC-AUC $\uparrow$ | ROC-AUC $\uparrow$ |
| GIN | $0.168 \pm 0.007$ | $77.9 \pm 1.0$ | $74.6 \pm 3.2$ | $66.0 \pm 2.1$ | $56.6 \pm 1.0$ |
| CAT+GIN | $\mathbf{0 . 1 0 1} \pm \mathbf{0 . 0 0 4}$ | $76.7 \pm 1.8$ | $\mathbf{7 9 . 5} \pm \mathbf{2 . 5}$ | $\mathbf{6 7 . 2} \pm \mathbf{1 . 8}$ | $\mathbf{5 8 . 2} \pm \mathbf{0 . 9}$ |
| GCN | $0.184 \pm 0.013$ | $76.7 \pm 1.4$ | $77.9 \pm 1.7$ | $66.1 \pm 2.4$ | $56.7 \pm 1.5$ |
| CAT+GCN | $\mathbf{0 . 1 2 3} \pm \mathbf{0 . 0 0 8}$ | $\mathbf{7 7 . 1} \pm \mathbf{1 . 6}$ | $\mathbf{7 9 . 2} \pm \mathbf{1 . 5}$ | $\mathbf{6 8 . 3} \pm \mathbf{1 . 7}$ | $\mathbf{5 7 . 9} \pm \mathbf{1 . 8}$ |
| GAT | $0.375 \pm 0.013$ | $76.6 \pm 2.0$ | $81.7 \pm 2.3$ | $66.2 \pm 1.4$ | $58.4 \pm 1.0$ |
| CAT+GAT | $\mathbf{0 . 2 0 1} \pm \mathbf{0 . 0 2 2}$ | $75.3 \pm 1.6$ | $79.3 \pm 1.6$ | $66.0 \pm 1.9$ | $58.3 \pm 1.3$ |
| Dataset $\rightarrow$ | MOLESOL | MOLTOXCAST | MOLLIPO | MOLTOX21 |  |
| $\downarrow$ Model | RMSE $\downarrow$ | ROC-AUC $\uparrow$ | RMSE $\downarrow$ | ROC-AUC $\uparrow$ |  |
| GIN | $1.105 \pm 0.077$ | $65.3 \pm 0.6$ | $0.717 \pm 0.016$ | $75.8 \pm 0.7$ |  |
| CAT+GIN | $\mathbf{0 . 9 8 5} \pm \mathbf{0 . 0 5 5}$ | $\mathbf{6 5 . 6} \pm \mathbf{0 . 5}$ | $0.798 \pm 0.031$ | $74.8 \pm 1.0$ |  |
| GCN | $1.053 \pm 0.087$ | $64.4 \pm 0.4$ | $0.748 \pm 0.018$ | $76.4 \pm 0.3$ |  |
| CAT+GCN | $\mathbf{1 . 0 0 6} \pm \mathbf{0 . 0 3 6}$ | $\mathbf{6 6 . 2} \pm \mathbf{0 . 8}$ | $0.771 \pm 0.023$ | $74.9 \pm 0.8$ |  |
| GAT | $1.037 \pm 0.063$ | $63.8 \pm 0.8$ | $0.728 \pm 0.024$ | $76.3 \pm 0.6$ |  |
| CAT+GAT | $1.09 \pm 0.048$ | $\mathbf{6 4 . 5} \pm \mathbf{0 . 8}$ | $0.754 \pm 0.021$ | $75.4 \pm 0.7$ |  |

## Strengths

- Better time complexity than other maximally expressive architecture for outerplanar graphs:
- 3-GNN (linear vs cubic)
- PlanE (linear vs quadratic)
- Very strong results in some datasets
- CAT can be applied to non-outerplanar graphs in linear time (without same guarantees)
- Recent work indicates CAT increases connectivity on the graph
- Most pharmaceutical molecules can be represented as outerplanar graphs.


## Weaknesses

- Experimental results not consistent: sometimes is even outperformed by base model on G (especially common for GAT). Considerations:
- CAT transformation introduces new "virtual" nodes and edges, so we have:
- longer dependencies
- GNN has to learn different representation for "virtual" nodes and edges
- No SOTA models implemented
- Guarantees restricted to outerplanar graphs (not very impactful)
- Would be good to generalize it to planar graphs (isomorphism still verifiable in polynomial time) like PlanE.


## QUESTIONS?

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