## Chapter 3

## **Specification models:**

High-level modeling techniques and related analysis methods for the computer assisted verification of DES



Discrete Event Systems Fall 2008

Computer Engineering and Networks Laboratory

## Motivation (1)

- Modern systems consist of many different components (HW + • SW !), yielding a very high degree of complexity
- Systems have to fulfill a set of requirements, defined by a ۲ specification as given by a contractor, customer or legislation
- What are these kinds of requirements? ۲
  - Functionality: Coke vending machines either delivers drink or returns my money
  - Performance: Voice-of-IP requires max. delay of a IP-packet < xxx msec.
  - Energy-consumption, heating characteristics
  - Reliability: 99,999% of emergency calls must be routed correctly
  - Safety / Security requirements (can been seen as part of functionality)
  - Economic requirements: costs, amortization etc.



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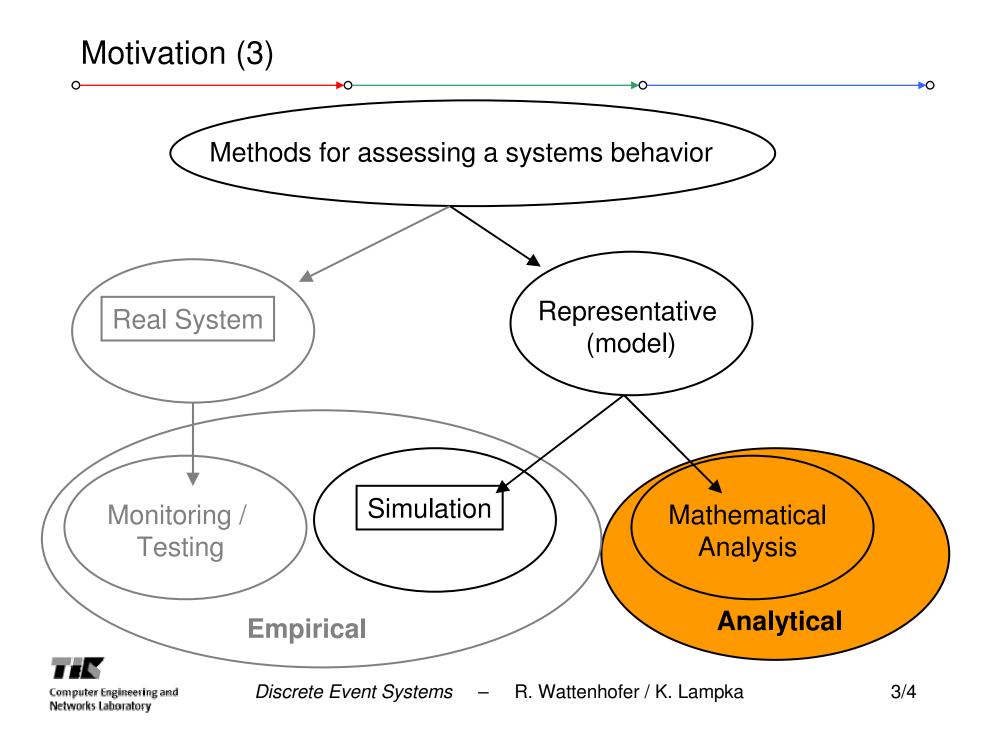
# Why do we need more sophisticated methods other than finite state machines?

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- States are atomic
  - no hierarchic structuring possible
  - usage of variables
- Partitioning of systems in (parallel operating) components not possible
  - modularization?
- FSM are easily very large and thus not human readable anymore



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Computer-assisted analysis/verification (1)

What's the general strategy for formally and automatically analyzing models?

- 1. Take a design of the system behavior, given as some high-level model such as an SDL specification, PN, network of TA, or .....
- 2. Take (a set of) formal system requirements, e.g. given in terms of a set of Message Sequence Charts, as safety or progress properties, ...
- 3. Validate that each requirement is indeed satisfied by the system design.



Computer-assisted analysis/verification (2)

### What are the techniques for formally analyzing models?

- 1. Theorem proving
  - Strategy: generate a formal proof that D satisfies  $\phi$ .
  - Applicable if design *D* can be represented in some adequate mathematical theory

### 2. State space exploration

- Strategy: check systematically and exhaustively each reachable state in D satisfies  $\phi$ .
- Applicable if the behavior of *D* can be finitely represented.
- One is enabled to show the presence and absence of errors!
- 3. Simulation or Testing of models
  - Strategy: Check whether  $\phi$  holds on some executions of *D*.
  - Applicable if *D* is in some sense executable.
  - One may be able to show the presence of errors, but not the absence!



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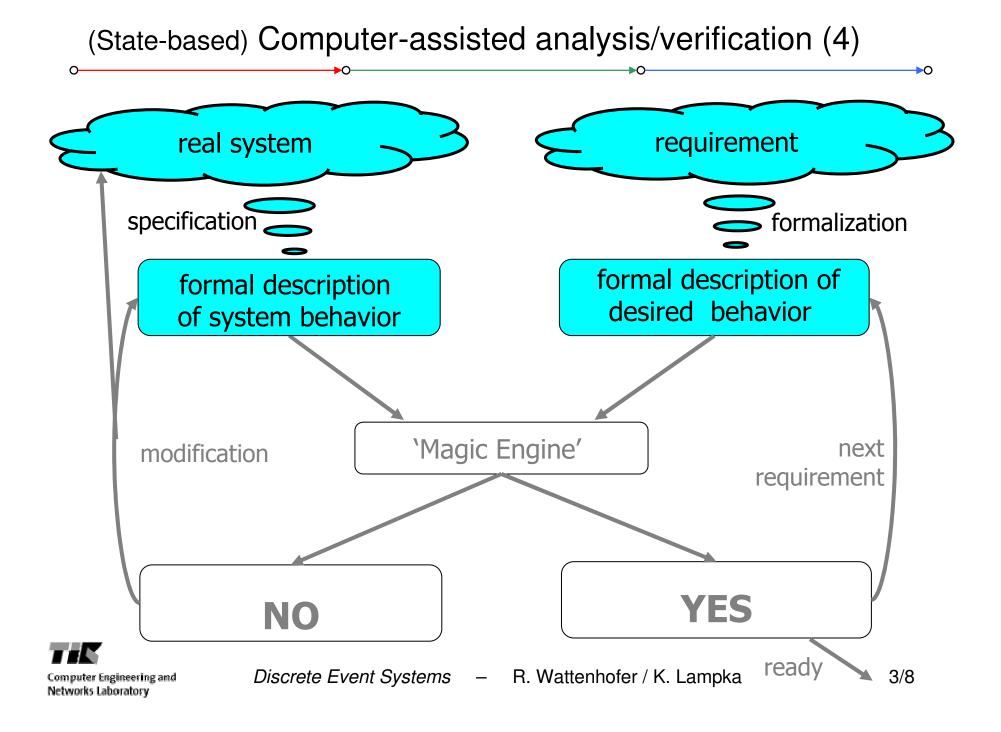
What are the requirements to be verified

(by state-based methods)?

- 1. Safety: A safety property to be verified asserts that a system under analysis never reaches a (set of) dedicated state, e.g. like error states, or in particular a deadlock. The mutual exclusion property is one of the most prominent examples of a safety property. (constraint on finite behaviour)
- 2. Liveness or progress: A liveness property guarantees that a system under analysis is executing a (set) of dedicated activities infinitely often (constraint on infinite behaviour)
- **3. Starvation** exists if there is an infinite run, where a dedicated action is never executed.

Dijkstra'71: Dining Philosophers Problem en.wikipedia.org/wiki/Dining\_philosophers\_problem



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(State-based) Computer-assisted analysis/verification (5)

Does a design *D* satisfy a requirement  $\phi$ ?

#### What are the practical obstacle?

- State Space Explosion: The number of possible state combinations *exponentially* in the number of concurrent processes or independent activities.
- Captures your model the reality ?
  - You can only assert those properties which are captured by the model.
  - Consistency check between model and reality ( = model validation, not covert here).

#### What's the principal obstacle?

• Decidability:

full generality it is undecidable whether D satisfies  $\phi$ . This depends on the specification, and the requirement.

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### Halting problem

#### Validation of model

Space and CPU-time is limited

## Validation ↔ Verification

- Validation ("doing right things") ٠
  - Comparing theory to reality
  - Make observation for providing evidence that the theory is correct, e.g.
    - show the absence of a specific error in different runs
    - show that program delivers correct result with respect to a given input
  - If possible try to find counter-examples (Falsification)
- Verification ("doing things right") ٠
  - Proof correctness of system design
  - formal model with unique interpretation
  - formal system, i.e. formal model and unique set of deductive rules for operating on (or transforming) the model.
  - If incorrect behavior is detected, counter-example is automatically provided



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## Specification Models (at glance)

- Systems have to fulfill a set of requirements, defined by a specification as given by a contractor, customer or legislation
- Showing (proving) a dedicated behavior of a system needs exhaustive analysis:
  - Testing, monitoring and simulation of real system or a model is in principle not sufficient (rare events?)
- Verification of systems require their formal description
- What does formal mean?
  - Model posses a unique interpretation (unique set of deductive <sup>Early</sup> versions of rules to operate on the model)
    UML-state
- Drawback: Reality far from trivial
  - => level of detail bounded by capabilities of (mathematically) handling a model (run-time and memory is limited)
- Abstraction:

Keep models simple as possible but as complex as necessary! You can only check what you have modeled!



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What are we doing?

- Lecture 23.10:
  - Specification and Description Language SDL

- Message Sequence Charts
- Related Analysis methods: The TAU-Tool-suite
- Lecture 30.10:
  - Petri Nets
  - Symbolic Analysis methods (for finite models)
- Lecture: 6.11:
  - Timed Automata
  - Introduction to model checking



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