

# The Internet Computer An Overview

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# • What is the Internet Computer? Consensus on the Internet Computer The Internet Computer Today







# What is the Internet Computer?



# What is the Internet Computer?

# Platform to run any computation, using blockchain technology for decentralisation and security



## **Coordination of nodes in** independent datacenters, jointly performing any computation for anyone

## ICP creates the Internet Computer blockchains

Guarantees safety and liveness of smart contract execution despite Byzantine participants





### ICP protocol

IP / Internet

Data Centers

\*\*\*\*\*\*\*

A REAL PROPERTY.

1 N N N

\*\*\*\*\*\*\* \*\*\*\*\*\*\*\*\*\*



## Canister smart contract

# **Canister Smart Contracts: Combination of Data and Code**





Code: WebAssembly bytecode



# **Developers and users interact directly with Canisters on the IC**

# DEPLOY



### Internet Computer



## Public cyberspace

# **Developers and users interact directly with Canisters on the IC**



### Internet Computer

### Nodes are partitioned into subnets

Canister smart contracts are assigned to different subnets

# Scalability: Nodes and Subnets



## Nodes are partitioned into subnets

Canister smart contracts are assigned to different subnets

## One subnet is special: it host the Network Nervous System (NNS) canisters which govern the IC

- ICP token holders vote on
- Creation of new subnets
- Upgrades to new protocol version
- Replacement of nodes
- . . .

# Scalability: Nodes and Subnets

![](_page_9_Picture_11.jpeg)

## Public key of NNS never changes, nodes in NNS share private key

## NNS generates key of subnets and certifies them

 Node in subnets use these keys to secure communication

# Chain Key Technology

![](_page_10_Picture_5.jpeg)

![](_page_11_Picture_0.jpeg)

# State:

canisters and their queues

## Inputs:

- new canisters to be installed,
- messages from users and other canisters

## **Outputs**:

 responses to users and other canisters

## **Transition function:**

- message routing and scheduling
- canister code

# **Each Subnet is a Replicated State Machine**

![](_page_11_Picture_16.jpeg)

![](_page_11_Picture_17.jpeg)

![](_page_12_Figure_1.jpeg)

# The Layers of the Internet Computer Protocol

### **Deterministic computation**

#### Message acquisition and ordering

![](_page_13_Picture_0.jpeg)

# **Consensus on the Internet Computer**

![](_page_13_Picture_2.jpeg)

## **Consensus Orders Messages**

![](_page_14_Picture_16.jpeg)

## Replicas may receive input messages in different orders, but must process them in the same order, for example 123466

# **Consensus Orders Messages**

## • Message (user $\rightarrow$ canister) • Message (canister $\rightarrow$ canister)

![](_page_15_Picture_3.jpeg)

![](_page_15_Picture_4.jpeg)

- Validity: all agreed upon blocks are valid

# **Consensus Properties**

Messages are placed in **blocks**. We reach agreement using a blockchain.

![](_page_16_Picture_8.jpeg)

Block x+1 Block x+2 Block x

The following properties must hold even if up to f < n/3 nodes misbehave

upon, they must have the same block *i*-th block is agreed upon

- Safety: For any *i*, If two (honest) replicas think that the *i*-th block is agreed
- Liveness: For any *i*, at some point every (honest) replica will think that the

![](_page_16_Picture_16.jpeg)

Neusen

![](_page_17_Picture_1.jpeg)

## • Message (user $\rightarrow$ canister) • Message (canister $\rightarrow$ canister)

## **Note:** We need more than one block maker in each round, otherwise the IC would not be fault tolerant!

## **Block Maker**

# A block maker selects available messages and combines them into a

30

## The notarization process ensures that a *valid* block proposal is published for every round

#### Step 1

Replica 1 receives a block proposal for height 30, building on some notarized height 29 block

![](_page_18_Figure_3.jpeg)

## Notarization

#### Step 2

Replica 1 sees that the block is valid, signs it, and broadcasts its notarization share

![](_page_18_Figure_7.jpeg)

![](_page_18_Picture_8.jpeg)

#### Step 3

Replica 1 sees that replicas 3 and 4 also published their notarization shares on the block

![](_page_18_Figure_11.jpeg)

#### Step 4

3 notarization shares are sufficient approval: the shares are aggregated into a single full notarization. Block 30 is now notarized, and notaries wait for height 31 blocks

![](_page_18_Figure_14.jpeg)

## Replicas may notary-sign multiple blocks to ensure that at least one block becomes fully notarized

#### Step 1

Replica 1 receives a block proposal for height 30, building on some notarized height 29 block

![](_page_19_Figure_3.jpeg)

## Notarization

#### Step 2

Replica 1 sees that the block is valid, signs it, and broadcasts its *notarization* share

![](_page_19_Figure_7.jpeg)

#### Step 3

Replicas 1 sees another height 30 block, which is also valid, and it broadcasts another notarization share

![](_page_19_Figure_10.jpeg)

#### Step 4

Both height 30 blocks get enough support to become notarized

![](_page_19_Figure_13.jpeg)

## Notarization

## Multiple notarized blocks may exist at the same height

![](_page_20_Figure_16.jpeg)

![](_page_20_Picture_17.jpeg)

![](_page_20_Figure_19.jpeg)

## At every height, there is a Random Beacon, an unpredictable random value shared by the replicas

#### Step 1

Replica 1 has Random Beacon 29 and wants to help constructing Random Beacon 30

![](_page_21_Picture_3.jpeg)

## Random Beacon

#### Step 2

Replica 1 signs RB29 using a threshold signature scheme, yielding a share of random beacon 30

![](_page_21_Picture_7.jpeg)

![](_page_21_Picture_8.jpeg)

#### Step 3

Replicas 1 sees that replica 2 also published a share of Random Beacon 30

![](_page_21_Figure_11.jpeg)

#### Step 4

2 random beacon shares are sufficient to reconstruct a full threshold signature, which is Random Beacon 30

![](_page_22_Figure_0.jpeg)

# **Block Maker Ranking**

The Random Beacon is used to rank block makers

1Replica 4Replica 24Replica 3Replica 33Replica 1Replica 42Replica 2Replica 1		RB 24	RB 25
<ul> <li>4 Replica 3 Replica 3</li> <li>3 Replica 1 Replica 4</li> <li>2 Replica 2 Replica 1</li> </ul>	1	Replica 4	Replica 2
3Replica 1Replica 42Replica 2Replica 1	4	Replica 3	Replica 3
2 Replica 2 Replica 1	3	Replica 1	Replica 4
	2	Replica 2	Replica 1

Round 25 Round 26

![](_page_22_Figure_5.jpeg)

#### Round 27

Round 28

![](_page_23_Picture_0.jpeg)

## Rounds are divided into time slots defining when block maker proposals are considered

![](_page_23_Figure_3.jpeg)

# **Notarization with Block Maker Ranking**

![](_page_23_Picture_5.jpeg)

![](_page_24_Picture_0.jpeg)

The block ranks can reduce the number of notarized blocks

#### Step 1

Replica 1 receives a rank-1 block proposal for height 30, building on some notarized height 29 block

![](_page_24_Figure_4.jpeg)

# **Notarization with Block Maker Ranking**

#### Step 2

Replica 1 is still in time slot 0, so not willing to notary-sign a rank-1 block yet

![](_page_24_Figure_8.jpeg)

#### Step 3

Replicas 1 sees a valid rank-0 height 30 block, and it broadcasts a notarization share

![](_page_24_Figure_12.jpeg)

#### Step 4

#### Eventually, only the rank 0 block becomes notarized

![](_page_24_Figure_16.jpeg)

![](_page_25_Picture_0.jpeg)

# **Notarization with Block Maker Ranking**

One notarized block *b* at a height *h* = Agreement up to *h* 

![](_page_25_Figure_8.jpeg)

How can we detect this...?

![](_page_26_Picture_0.jpeg)

# **Notarization with Block Maker Ranking**

## Synchronous communication $\rightarrow$ Forks can be removed

![](_page_26_Figure_6.jpeg)

![](_page_26_Figure_8.jpeg)

#### 37 Rank 0

![](_page_27_Picture_0.jpeg)

## Partially synchronous communication $\rightarrow$ Forks cannot be removed!

![](_page_27_Picture_2.jpeg)

![](_page_27_Picture_3.jpeg)

# **Notarization with Block Maker Ranking**

## Replicas create finalization shares if they did not sign any other block at that height

![](_page_28_Figure_1.jpeg)

Replica 1 observes that block *b* is fully notarized and will no longer notary-sign blocks at height  $\leq$  30

# Finalization

#### Step 2

![](_page_28_Figure_5.jpeg)

### Step 3

Since replica 1 did not notary-sign any other block than block b, it signs block *b*, creating a finalization-share on *b* 

![](_page_28_Figure_8.jpeg)

#### Replica 1 did not notary-sign any height 30 block other than b

#### Step 4

Replicas 2 and 4 also cast finalization shares on block *b* 

![](_page_28_Picture_12.jpeg)

#### Step 5

3 finalization-shares are sufficient approval: the shares are aggregated into a single full finalization

![](_page_28_Figure_15.jpeg)

## Finalization on block b at height h = Proof that no other block is notarized at height h

![](_page_29_Picture_1.jpeg)

![](_page_29_Picture_2.jpeg)

## **Proof**:

- 1. A full finalization on *b* requires *n-f* replicas to finality-sign (by construction)
- - that  $\leq f$  replicas are corrupt)
- h (by construction)
- 5. A full notarization requires *n-f* notarization-shares (by construction)
- sufficient to reach the notarization threshold of *n-f* (by 4. & 5.)

# Safety of Finalization

If block b at height h is finalized, then there is no finalized block  $b' \neq b$  at height h.

2. At least *n-2f* of the *n-f* replicas that finality-signed *b* must be honest (by assumption

3. An honest replica that finality-signed b did not notary-sign any other block at height

4. At least *n-2f* replicas did not notary-sign any height *h* block other than *b* (by 2. & 3.) 6. The *n-(n-2f) < n-f* remaining replicas that may have notary-signed a block b' are not

![](_page_31_Picture_0.jpeg)

# The Internet Computer Today

![](_page_31_Picture_6.jpeg)

![](_page_32_Picture_0.jpeg)

# Live Since May 2021!

# Currently 375 machines by 53 node providers

Network Status Operational \*

![](_page_32_Figure_11.jpeg)

### https://dashboard.internetcomputer.org/

- Disseminating messages among all nodes in the same subset • Exchanging canister and control messages between subnets Scheduling and concurrent execution of canister messages Catching up after a node has been offline for a while Handling churn (adding and removing nodes) • Guaranteeing consistency (different users need a consistent view of data and operations) Upgrading to next protocol version

- Creating new subnets
- Load balancing

# Many distributed systems problems

![](_page_34_Picture_0.jpeg)

# Fast Growing Blockchain Ecosystem

#### Average block time:

Finality:

#### TX per second:

Validation data:

## Internet Computer vs. ...

![](_page_35_Picture_5.jpeg)

#### 1 block / 10 minutes

#### 1 hour

7

380 GB

![](_page_35_Picture_10.jpeg)

![](_page_35_Picture_11.jpeg)

#### 1 block / 15 seconds

#### 3 minutes

15

#### 550 GB

![](_page_35_Picture_17.jpeg)

### 30 blocks / second

#### 1-3 seconds

#### 11,500 (write) / 250,000 (read)

#### 48 bytes

## Infographic: <u>here</u>

#### Technical Library: <u>here</u> (videos of talks) and <u>here</u> (blogposts) lacksquare

## 200,000,000 CHF Developer Grant Program <u>here</u>

### • DFINITY SDK: <u>here</u>

# **More information**

![](_page_36_Picture_5.jpeg)

![](_page_36_Picture_6.jpeg)

INTERNET COMPUTER		DFINITY.org Forum Support 🎔 🕓 🔍	Search	
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For: First-Timers			ON THIS PAGE	
For: Developers		The Internet Computer is a blockchain that runs at web speed with	For: First-Timer	s
For: Protocol Enthusiasts		unbounded capacity.	For: Developers	
		As a crash course in blockchain history, Bitcoin created digital gold.	For: Protocol	ack
Quick Start	>	Then, in the next step of the evolution, Ethereum developed smart	Enthusiasts	edb
Example Code	>	contracts and pioneered DeFi and NFT use cases.		Ъ.
Developer Docs	>	The Internet Computer is the third major blockchain innovation — a		
Protocol Docs	\$	blockchain that scales smart contract computation, runs them at		
FIGUEULDUCS		web speed, processes and stores data efficiently, and provides		
General Docs	>	powerful software frameworks to developers. By making this		
Additional Resources	>	possible, the Internet Computer enables the complete reimagination		
		of now systems and apps operate.		

![](_page_36_Picture_8.jpeg)

#### Introducing the Internet **Computer Interface** Specification

It details how services and users communicate through the Internet Computer, and enables the community to for internet-scale services. create new development tools.

Mar 19 · 4 min read

![](_page_36_Picture_12.jpeg)

	102 GRANTEES
Grantee	
AEDILE	
Project management dapp	
AGRYO	
Global risk intelligence for agriculture	
ASTROX	
Dart developer tools and "mini apps" framework	
B9 LABS	
Developer onboarding documentation	
BAUCTION	
Decentralized and transparent auction platform	

![](_page_36_Picture_14.jpeg)

#### A Closer Look at Software **Canisters, an Evolution of** Smart Contracts

Canisters are smart contracts that scale — interoperable compute units designed

![](_page_36_Picture_17.jpeg)

![](_page_36_Picture_18.jpeg)

![](_page_36_Picture_19.jpeg)

#### A Technical Overview of the Internet Computer

An explanation of the blockchain network's infrastructure, and how canister smart contracts enable web services to scale without bound.

Sep 18, 2020 · 12 min read

![](_page_36_Picture_23.jpeg)

![](_page_37_Picture_1.jpeg)