

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich

Distributed Computing



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Computer Systems

Assignment 6

Assigned on: November 2, 2018

1	Clock	Sync

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1.1 Clock Synchronization

- a) Assume you run NTP to synchronize speakers in a soccer stadium. Each speaker has a radio downlink to receive digital audio data. However, there is no uplink! You decide to use an acoustic signal transmit by the speaker. To synchronize its clock, a speaker first plays back an acoustic signal. This signal is picked up by the NTP server which responds via radio. The speaker measures the exact time that passes between audio playback and radio downlink response. What is likely the largest source of error?
- b) What are strategies to reduce the effect of this error source?
- c) Prove or disprove the following statement: If the average local skew is smaller than x, then so is the average global skew.
- d) Prove or disprove the following statement: If the average global skew is smaller than x, then so is the average local skew.

Basic	

1.2 Time Difference of Arrival

Assume you are located on a line y=-x+8 km in the two dimensional plane. You receive the GPS signals from satellites A and B. Both signals are transmitted exactly at the same time t by both satellites. You receive the signal from satellite A 3.3 µs before the signal of satellite B. At time t, satellite A is located at $p_A = (6 \text{ km}, 6 \text{ km})$ and satellite B is located at $p_B = (2 \text{ km}, 1 \text{ km})$, in the plane.

- a) Formulate the least squares problem to find your location.
- b) Are you more likely to be at position (2 km, 6 km) or (4 km, 4 km)?
- c) What is the time when receiving the signal from satellite B?

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1.3 Clock Synchronization: Spanning Tree

Common clock synchronization algorithms (e.g. TPSN, FTSP) rely on a spanning tree to perform clock synchronization. Finding a good spanning tree for clock synchronization is not trivial. Nodes which are neighbors in the network graph should also be close-by in the resulting tree. Show that in a grid of $n = m \times m$ nodes there exists at least a pair of nodes with a stretch of at least m. The stretch is defined as the hop distance in the tree divided by the distance in the grid.

1.4 NTP Programming

Write a Linux program that prints the current UTC time and the maximum error.

Hint: Have a look at the manpage for adjtimex.

2 Consistency and Logical Clocks

Quiz _

2.1 Different Consistencies

Prove or disprove the following statements:

- a) Neither sequential consistency nor quiescent consistency imply linearizability.
- b) If a system has sequential consistency and quiescent consistency, it is linearizable.

Basic _

2.2 Measure of Concurrency from Vector Clocks

You are given two nodes that each have a vector logical clock that additionally logs the clock state upon receiving a message (see Algorithm 1).

Algorithm 1 Vector clocks with logging

- 1: (Code for node u)
- 2: Initialize $c_u[v] := 0$ for all other nodes v.
- 3: Upon local operation: Increment current local time $c_u[u] := c_u[u] + 1$.
- 4: Upon send operation: Increment $c_u[u] := c_u[u] + 1$ and include the whole vector c_u as d in message.
- 5: Upon receive operation: Extract vector d from message and update $c_u[v] := \max(d[v], c_u[v])$ for all entries v. Increment $c_u[u] := c_u[u] + 1$. Save the vector c_u to the log file of node u.

Assume that exactly one message gets send from one to the other node. Given the logs and current vector states of both nodes, write a short program that calculates the measure of concurrency as defined in the script (Definition 3.30). You can use your favorite programming language. The example solution will be in Python.

Advanced _

Generalize your program to any number of messages exchanged between the nodes.