





 $\mathrm{HS}\ 2011$

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Discrete Event Systems

Exercise Sheet 5

1 Counter Automaton

A push-down automaton is basically a finite automaton augmented by a stack. Consider a finite automaton that (instead of a stack) has an additional *counter* C, i.e., a register that can hold a single integer of arbitrary size. Initially, C = 0. We call such an automaton a *Counter Automaton* M. M can only increment or decrement the counter, and test it for 0. Since theoretically, all possible data can be coded into one single integer, a counter automaton has unbounded memory. Further, let L_{count} be the set of languages recognized by counter automata.

- a) Let L_{reg} be the set of regular languages. Prove that $L_{reg} \subseteq L_{count}$.
- b) Prove that the opposite is not true, that is, $L_{count} \not\subseteq L_{reg}$. Do so by giving a language which is in L_{count} , but not in L_{reg} . Characterize (with words) the kind of languages a counter automaton can recognize, but a finite automaton cannot.
- c) Which automaton is stronger? A counter automaton or a push-down automaton? Explain your decision.

2 Tandem Pumping

For the following languages, determine whether they are context free or not. Prove your claims!

- a) $L = \{a^i b^j c^k \mid 0 < i < j < k\}$
- b) $L = \{x \mid x \in \{0,1\}^*, \text{ and } x \text{ contains an even number of '0' and an even number of '1'}\}$

3 Context Free Grammars

Are the following languages context-free? If so, give a CFG describing L.

- a) $L = \{w \# x \# y \# z \mid w, x, y, z \in \{a, b\}^* \text{ and } |w| = |z|, |x| = |y|\}$
- **b)** $L = \{w \# x \# y \# z \mid w, x, y, z \in \{a, b\}^* \text{ and } |w| = |y|, |x| = |z|\}$

4 Push Down Automata

For each of the following context free languages, draw a PDA that accepts L.

- **a)** $L = \{a^i b^j a^j b^i \mid i, j > 0\}$
- **b)** $L = \{u \mid u \in \{0, 1\}^* \text{ and } u^{reverse} = u\} = \{u \mid ``u \text{ is a palindrome}"\}$
- c) $L = \{u \mid u \in \{0, 1\}^* \text{ and } u^{reverse} \neq u\} = \{u \mid ``u \text{ is no palindrome''}\}$

5 Designing Turing Machines

Alice is very happy because she was accepted for an internship at Tintel, one of the world's leading processor manufacturers. Unfortunately, she has only attended the famous DES lecture during her studies at ETH and knows nothing about electronic circuits. Therefore, she wants to solve her first assignment using a Turing Machine – please assist her:

Alice is asked to implement a *binary to unary converter*. This converter takes a number a in binary (alphabet $\{0,1\}$) and converts it to a unary number u (alphabet $\{1\}$). Initially, the TM head points to the MSB of a. At the end, the head should point to the right-most digit of u.

Provide a plain text description of your TM as well as a finite state machine controlling the tape head. Use the following notation for transitions:

- $\alpha \rightarrow \beta \mid \gamma'$ read α from the tape at the current position, then write a β and finally move left if $\gamma = L$ or move right if $\gamma = R$.
- ' $\alpha \mid \gamma$ ' abbreviation for transitions of the form $\alpha \to \alpha \mid \gamma$ (these transitions do not modify the content of the tape).

Hint: The number n in unary representation consists of n ones. Also, you might want to extend the alphabet Γ to put temporary symbols on the tape.