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## **Discrete Event Systems**

Exercise Sheet 3

## Finite Automata and Regular Languages [Exam] 1

Consider the NFA A in Figure 1 and further assume  $\Sigma = \{0, 1\}$ .



Figure 1: NFA A.

- a) Transform the NFA into an equivalent DFA using the powerset construction of the lecture. (*Hint:* Only construct states which are necessary!)
- **b)** Which regular language is accepted by the automaton A?

## 2 Pumping Lemma [Exam]

Are the following languages regular? Prove your claims!

- a)  $L_1 = \{0^a 1^b 0^c 1^d \mid a, b, c, d \ge 0 \text{ and } a = 1, b = 2 \text{ and } c = d\}$
- **b)**  $L_2 = \{0^a 1^b 0^c 1^d \mid a, b, c, d \ge 0 \text{ and if } a = 1 \text{ and } b = 2 \text{ then } c = d\}$

## 3 Transforming Automata [Exam]

Consider the DFA in Figure 2 over the alphabet  $\Sigma = \{0, 1\}$  accepting the language L. Let  $\Phi(L)$  be defined as

$$\Phi(L) := \{ w \in \Sigma^* \mid \exists x \in \Sigma^*, |x| = |w| \text{ and } wx \in L \}$$

In other words,  $\Phi(L)$  is the set of the front halves of all words in L.



Figure 2: DFA B.

- a) Give a regular expression for the language L accepted by the automaton B. If you like, you can do this by ripping out states as presented in the lecture (slide 1/84 ff.).
- **b)** Construct a DFA which accepts a word w if and only if  $w \in \Phi(L)$ .