

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich Distributed Computing

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Ad Hoc And Sensor Networks Sample Solution to Exercise 7

Assigned: November 8, 2010 Due: November 15, 2010

1 Slotted Aloha

We define the function $P: \mathbb{R}^2 \to \mathbb{R}$ as

 $P(p,n) := \Pr \operatorname{success} = n \cdot p(1-p)^{n-1}.$

For a fixed p, P(p, n) is monotone increasing for $n \leq -1/\ln(1-p)$ and monotone decreasing for $n \geq -1/\ln(1-p)$ and therefore P(p, n) is minimized either at n = A or at n = B for $n \in [A, B]$. Therefore, we have to find

$$p_{\text{opt}} := \max_{p} \left(\min \left\{ P(p, A), P(p, B) \right\} \right).$$

For a fixed n, P(p, n) is monotone increasing for $p \leq 1/n$ and monotone decreasing for $p \geq 1/n$ (for $p \in [0, 1]$). Furthermore, $P(1/A, A) \geq P(1/A, B)$ and $P(1/B, B) \geq P(1/B, A)$ for $B \geq A + 1$ and therefore the intersection between P(p, A) and P(p, B) is between the maxima of P(p, A) and P(p, B), respectively. Thus p_{opt} is found where $P(p_{\text{opt}}, A) = P(p_{\text{opt}}, B)$.

$$A * p_{\text{opt}} * (1 - p_{\text{opt}})^{A-1} = B * p_{\text{opt}} * (1 - p_{\text{opt}})^{B-1}$$
$$\frac{A}{B} = (1 - p_{\text{opt}})^{B-1-(A-1)} = (1 - p_{\text{opt}})^{B-A}$$
$$p_{\text{opt}} = 1 - \sqrt[B-A]{\frac{A}{B}}.$$

For A = 100 and B = 200, we get

$$p_{\rm opt} = 0.006908 = \frac{1}{144.8}.$$

2 Broadcast

Student A is right.

An exemplary algorithm:

Source originating the broadcast: Transform the message m as follows: Replace a 1 with 10 and append 11 at the end and at the front of a message, i.e. message m = 10110 becomes message $m' = 11\ 10010100\ 11$. Transmit "Hello" in round i if bit i of m' is 1. If a node is not the source it waits until it detects twice a non-free channel for two consecutive rounds. It decodes a non-free channel as 1 and a free channel as 0. It can easily reconstruct the message m by ignoring 11 at the beginning and end and replacing 10 with 1 for the bits received in between the first received 11 and the second 11. As soon as a node decoded the entire message m, it starts to transmit the same m' in the same way as the source.