Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich



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Computer Engineering II

Exercise Sheet Chapter 11

Quiz _

1 Quiz

- a) If requesting a lock is very expensive, how would you implement a linked list set?
- b) What properties do (good) hash functions have? List as many as you can!
- c) How would you implement a hash map supporting inserting multiple values per key?
- d) Which of the implementations for a list-based set is FIFO fair?

Basic .

2 Livelock

In the lecture we discussed how to implement a Set using a linked list and the concept of optimistic synchronization. The main trick was to only lock affected parts of the list once a change should be applied. Are there bad situations in which the algorithm works badly?

- a) Is there a scenario in which two (or more) threads deadlock? If yes: give an example. If no: argue why.
- **b)** Is there a scenario in which one thread never succeeds in removing a node? If yes: give an example. If no: argue why.

3 Old Exam Question: Fine-Grained Locking

The goal of this exercise is to implement a heap with mutual exclusion. A heap is a binary tree, in which the value of the parent is smaller than the values of its children. The heap is stored in an array, with the root at index 1 and the children of a node i are $LEFT(i) = 2 \cdot i$ and $RIGHT(i) = 2 \cdot i + 1$. The basic functionality is implemented in Algorithm 1 and Algorithm 2.

Algorithm 1 Insert value	Algorithm 2 Remove smallest value
1: $i = 1$	1:
2:	2: $ret = A[1]$
3: while $A[i] = null do$	3: i=1
4:	4: $A[1] = \infty$
5: $next = smallestChild(i)$	5:
6:	6: while $A[i] := null do$
7: if $(A[i] > value)$ then	7:
8: exchange A[i] and value	8: $next = smallestChild(i)$
9: end if	9:
10:	10: if $(A[next] != null)$ then
11: $i = next$	11: exchange A[i] and A[next]
12:	12: else
13: end while	13: $A[i] = null // Mark as not used$
14:	14: end if
15: $A[i] = value$	15:
16:	16: $i = next$
	17:
	18: end while
	19:
	20: return ret

- a) (4 Points) How would you implement coarse-grained locking? What consequences does this have for concurrent access by multiple processes?
- b) (8 Points) Complete the skeleton of the code in Algorithm 1 and Algorithm 2 to implement hand-over-hand locking. You may use LOCK(j) and UNLOCK(j), which lock/unlock the jth element in the array. Not all lines are needed. You may use multiple statements per line.
- c) (5 Points) Is your implementation deadlock free? Argue why deadlocks are not possible or provide an example of a deadlock.
- d) (3 Points) When using hand-over-hand locking the root is always locked at the beginning of each operation. Could you use a different locking mechanism to avoid this contention of the root?